

**MA222 - Computational Linear Algebra
Problem Sheet - 3**

Block Matrices and Algorithms

1. Adapt strass so that it can handle square matrix multiplication of any order.
Hint: If the "current" A has odd dimension, append a zero row and column.
2. Prove that if

$$A = \begin{bmatrix} A_{11} & \cdots & A_{1r} \\ \vdots & \ddots & \vdots \\ A_{q1} & \cdots & A_{qr} \end{bmatrix}$$

is a blocking of the matrix A , then

$$A^T = \begin{bmatrix} A_{11}^T & \cdots & A_{q1}^T \\ \vdots & \ddots & \vdots \\ A_{1r}^T & \cdots & A_{qr}^T \end{bmatrix}.$$

3. Suppose n is even and define the following function from \mathbb{R}^n to \mathbb{R} :

$$f(x) = x(1 : 2 : n)^T x(2 : n) = \sum_{i=1}^{n/2} x_{2i-1} x_{2i}$$

- (a) Show that if $x, y \in \mathbb{R}^n$ then

$$x^T y = \sum_{i=1}^{n/2} (x_{2i-1} + y_{2i})(x_{2i} + y_{2i-1}) - f(x) - f(y)$$

- (b) Now consider the n -by- n matrix multiplication $C = AB$. Give an algorithm for computing this product that requires $n^3/2$ multiplies once f is applied to the rows of A and the columns of B .

4. Prove Lemma 1.3.2 for general s .

Hint. Set $p_T = p_1 + \cdots + p_{\gamma-1}$ $\gamma = 1 : s + 1$ and show that $c_{ij} = \sum_{\gamma=1}^s \sum_{k=p_{\gamma}+1}^{p_{\gamma+1}} a_{ik} b_{kj}$.

5. Use Lemmas 1.3.1 and 1.3.2 to prove Theorem 1.3.3. In particular, set

$$A_{\gamma} = \begin{bmatrix} A_{1\gamma} \\ \vdots \\ A_{q\gamma} \end{bmatrix} \quad \text{and} \quad B_{\gamma} = [B_{\gamma 1} \quad \cdots \quad B_{\gamma r}]$$

and note from Lemma 1.3.2 that

$$C = \sum_{\gamma=1}^s A_{\gamma} B_{\gamma}.$$

Now analyze each $A_{\gamma} B_{\gamma}$ with the help of Lemma 1.3.1.
